

(Verifiable) Delay Functions from Lucas Sequences

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Verifiable Delay Functions

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 - resource-efficient blockchains [CP19]
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• constructions:

- hidden order groups [Pic19, Wes20]
- SNARKs [Bon+18]
- isogenies [Fee+19, Sha19, CRT21]
- lattices [LM23, CLM23]
- squaring in a finite field [KMT22]

(Modular) Lucas Sequences

- $U(P, Q)$: $U_0 = 0, U_1 = 1, U_i = PU_{i-1} - QU_{i-2}$
- $V(P, Q)$: $V_0 = 2, V_1 = P, V_i = PV_{i-1} - QV_{i-2}$

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- primality tests, factorization [Ric 85, W:182]
- LUC cryptosystems [LS93, MN81, BBL85]

Construction

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let $\omega, \bar{\omega} \in \mathbb{Z}_N[\sqrt{P^2 - 4Q}]$ be roots of $x^2 - Px + Q$

then
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- $g_T(P, Q) = \omega^{2^T}$

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$$f_T: \mathbb{Z}_N \rightarrow \mathbb{Z}_N$$
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$$\frac{U_{2^T}^2 - DV_{2^T}^2}{4} = Q^{2^T} \quad [PB91]$$

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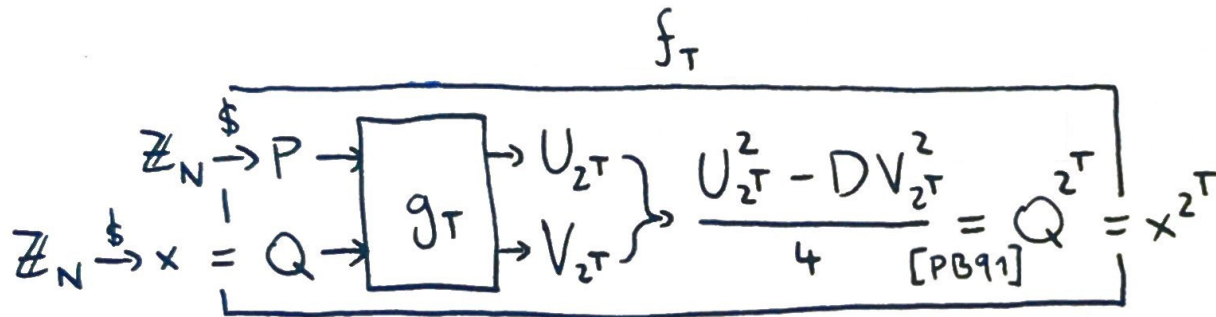
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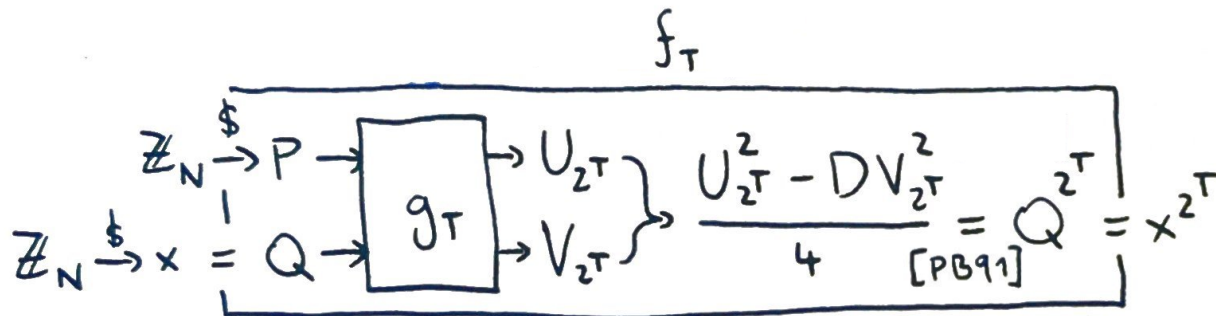
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• conjectured: $RSW < LCS$

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Pietrzak's protocol [Pie19]

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$$x \stackrel{T}{\rightarrow} y$$

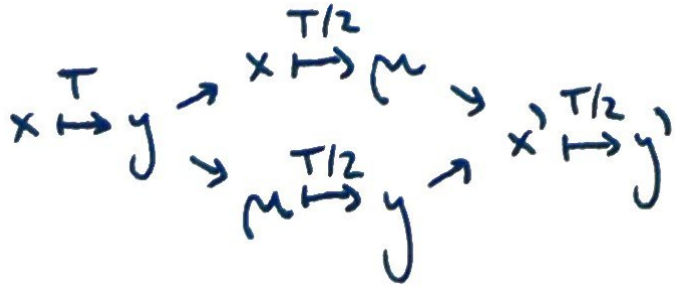
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$$\begin{array}{l} x \xrightarrow{T} y \rightarrow \\ \quad \quad \quad \rightarrow \begin{array}{l} x \xrightarrow{T/2} \mu \\ \mu \xrightarrow{T/2} y \end{array} \end{array}$$

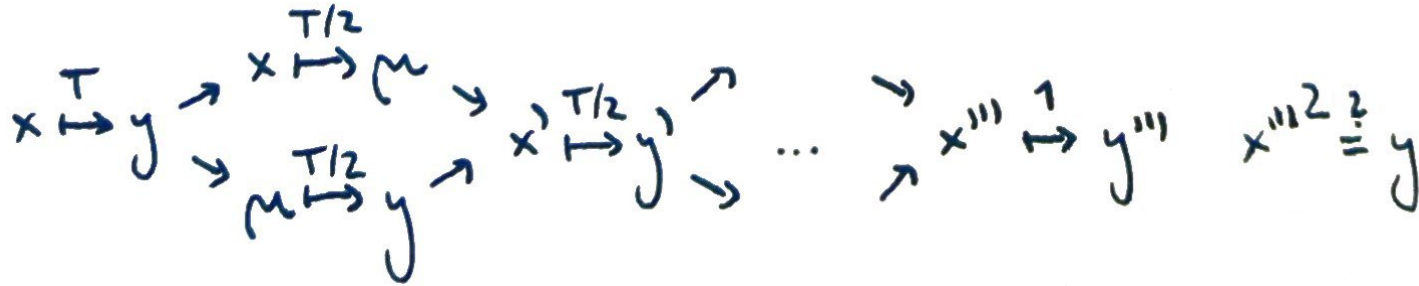
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$\mathbb{Z}_{p-1}^{(z)}(\cdot)$ is a prime
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strong primes

$\mathbb{Z}_{p-1}^{(a)}(\cdot)$ and $\mathbb{Z}_{p+1}^{(a)}(\cdot)$ have
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group
membership
verifiability

$\sqrt[z]{\mu}$ is a previous
member of the sequence

This work

$$\mathbb{Z}_N^{(a)}[\sqrt{D}](\cdot)$$

$$w \mapsto w^{z^T}$$

strong primes

$\mathbb{Z}_{p-1}^{(a)}(\cdot)$ and $\mathbb{Z}_{p+1}^{(a)}(\cdot)$ have
no subgroups of small order

$$a \in \mathcal{O}(1)$$

prover gets $\sqrt[z]{\mu}$ by
delaying cofactor clearing

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 - verifiability of primality of Mersenne numbers
 - non-primality of certificates for Proth's numbers [Hof+23]

$$2^{82,589,933} - 1$$

Mersenne

$$10223 \cdot 2^{31,172,165} + 1$$

Proth

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Thank you for your attention!

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